# The many facets of mathematical Optimization

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#### A plan?

What is optimization?

Is optimization useful?

Is optimization easy?

Is optimization always hard?

Can optimization help with your diet?

Can we trust optimization?

Can optimization beat computer chip makers?

# 1. What is Optimization?

#### Modelling → Decision

- 1. Model a real-word situation mathematically
- 2. Choose the **best** feasible decision

Decision → vector of variables

Best → objective fonction

Feasible → constraints

Variables+objective+constraints → optimisation A mathematical **tool** for decision making

#### Mathematical formulation

Variables:  $x \in \mathbb{R}^n$ 

Objective function:  $f: \mathbb{R}^n \mapsto \mathbb{R}: x \mapsto f(x)$ 

Constraints:  $D \subseteq \mathbb{R}^n$ 

 $\min_{x \in \mathbb{R}^n} f(x)$  such that  $x \in D$ 

#### A few simple examples

Find rectangular box of a given volume V with the minimum surface

$$min_{(a,b,c)\in R^3} 2ab + 2bc + 2ca \ s.t. \ abc = V$$

Find the diameter of a shape S in the plane

$$\max_{x \in R^2, y \in R^2} ||x - y|| \text{ s.t. } x \in S, y \in S$$

Find the fastest path to your destination (in your GPS device) – *binary variables* 

### Terminology

#### **Applied Mathematics**

- → Operations Research
  - = Quantitative methods for decision making
  - = Management Science
- = Industrial Engineering
  - Mathematical Programming
    - = Optimization = Optimisation
    - = Mathematical Optimization
    - = Optimalisation

#### Optimization: a vast domain

- Deterministic or stochastic problem
- Certain or uncertain data
- Single or multiple objective
- Explicit or black-box characteristics
- Continuous or discrete
- Linear or non-linear
- Constrained or unconstrained
- etc.

Presentation will focus on a small part

# 2. Is Optimization useful?

#### A wide variety of applications

- Planning and scheduling
   Production, timetabling, crew scheduling
- Design
   Structural design, network design
- Economics and finance
   Portfolio optimization, Nash equilibria
- Location and transport
   Facility location, routing, tours

#### INFORMS Edelman award 2008

#### **Dutch railways**

In 2006: Schedule optimisation for passengers, trains and personnel

#### After one year:

- Delays reduced by 17%
- Passengers increased by 10%
- Profit increased by € 40M

### Engineering applications

- Airbus: structural optimization of wings (min weight for given strength)
- Radiotherapy: dose optimization (modulate a moving beam)
- Spectrum optimization in telecommunications (your next modem)
- Astrophysics : optimize diffraction masks (planet detection)
- And much much more ...

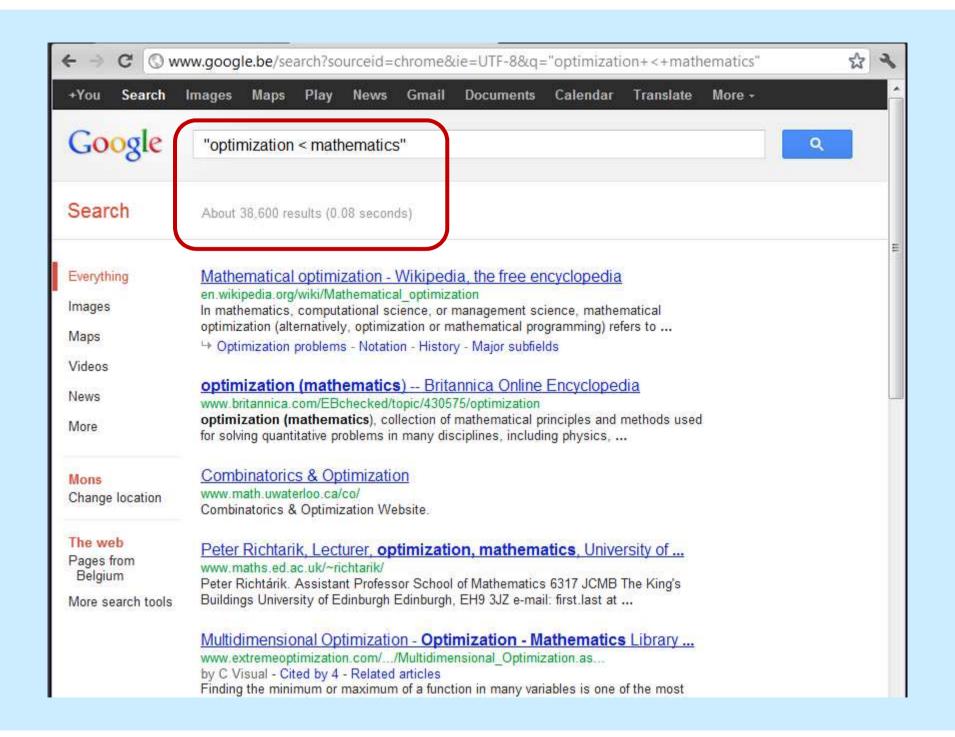
### Optimization is important

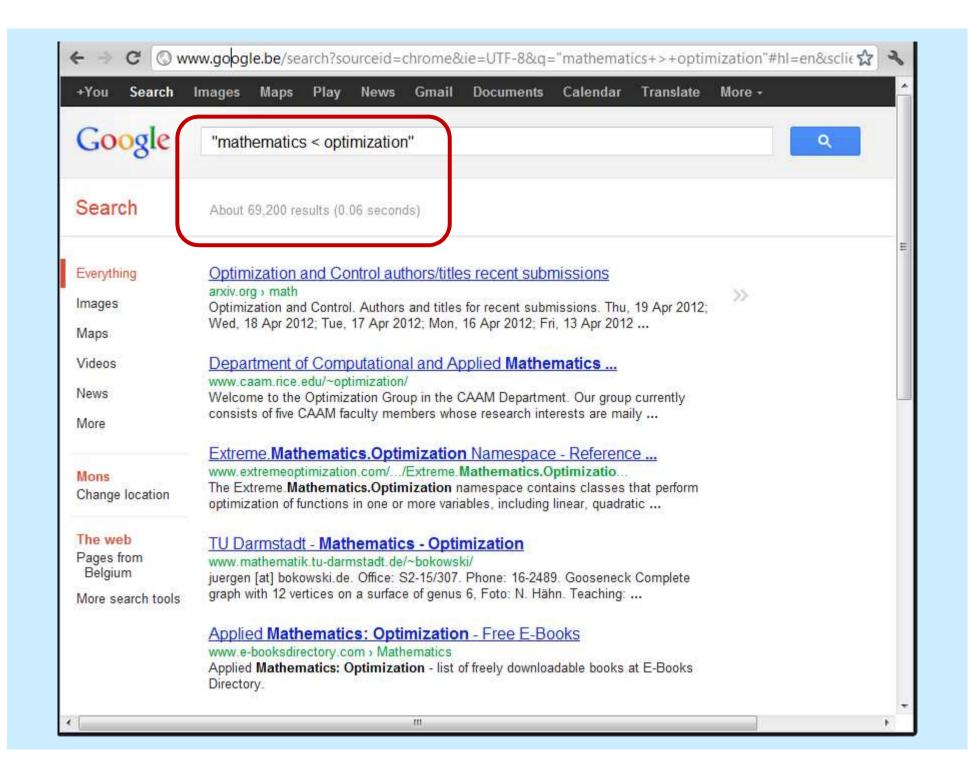
**Mathematics** 

Versus

**Optimization** 

Let us ask the modern Oracle ...





### Googlefight: verdict

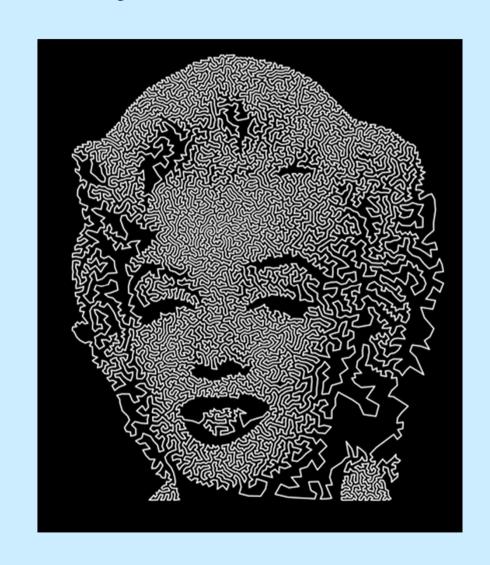
**Mathematics < Optimization : 69 200** 

versus

Optimization < Mathematics : 38 600

### Opt-Art, by R. Bosch

A single loop!



# 3. *Is Optimization easy?*

### Algorithmic black-box optimization

$$f^* = \min_{x \in \mathbb{R}^n} f(x)$$
 such that  $x \in D$ 

- Rules of the game: the black box
  - Give me points in  $D: x^{(1)}, x^{(2)}, x^{(3)}, ...$
  - I will return the values of  $f: f(x^{(1)}), f(x^{(2)}), \dots$
  - Nothing else is known about the function
- Goal: find a good solution quickly
  - As few evaluations as possible
  - With prescribed accuracy  $\varepsilon$ :  $f(x^{(n)}) f^* < \varepsilon$

#### Looking at the worst case

Typical theorem:

For a given class of functions, any method must take at least some number of evaluations on some function

Must hold for any method

Main tool for proof: resisting oracle

#### Some games cannot be won

Typical theorem:

For a given class of functions, any method must take at least some number of evaluations on some function

- For some functions, game is always lost
  - discontinuous functions
  - functions with very quick variations
  - tiny and complicated domain D

### Optimization is hard: the theorem

Among the class of functions with n variables defined on the unit box  $\{x \mid 0 \le x_i \le 1 \ \forall i\}$  whose variation is not too large in the sense

$$|f(x) - f(y)| < L ||x - y||_{\infty}$$

for any method there exists a function such that accuracy  $\varepsilon$  is not reached before

$$\left\lfloor \frac{L}{2\varepsilon} \right\rfloor^n$$
 evaluations

### Far-reaching consequences

 $\left\lfloor \frac{L}{2\varepsilon} \right\rfloor^n$  evaluations

Consider an easy problem with L=2 and  $\mathcal{E} = 1\%$ Some 1-variable problem *must* take 100 iter. Some 10-variable problem *must* take  $10^{20}$  iter. Life is too short for such problems!

- Uniform grid search is worst-case optimal
- Result does not really improve for smoother f

# 3. Is Optimization always hard?

#### Optimization is not always hard

No hope of guaranteeing resolution of optimization problems with 10 variables

Two ways to react:

Ignore theorem

Try you luck!

Nullify theorem

Force your luck! new problem class

#### A new class: linear optimization

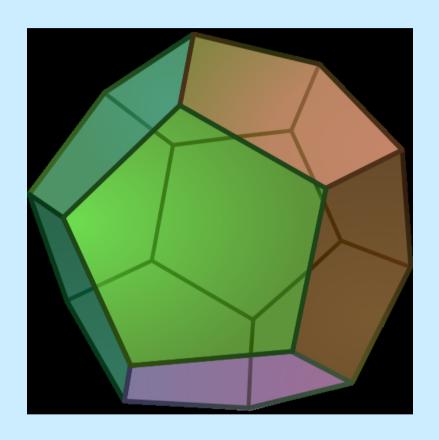
We voluntarily restrict ourselves to a problem class involving only linear functions

$$\min_{x \in \mathbb{R}^n} f(x)$$
 such that  $x \in D$ 

Linear objective function f

Linear constraints, which gives polyhedron D

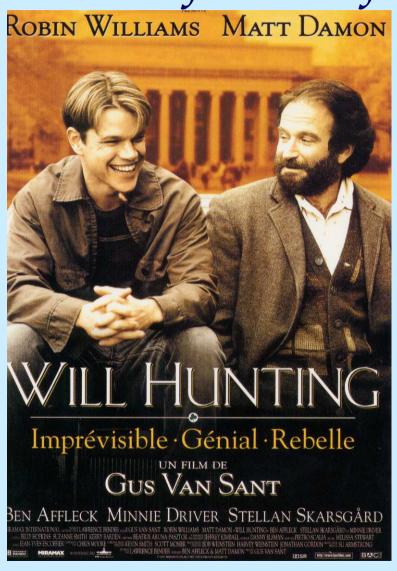
#### A polyhedron



Application examples:

Production problems, assignment problems, transport problems, flow problems, etc.

#### The father of linear optimization



Georges B. Dantzig (1914-2005) American statistician Stanford Professor

Originated the « difficult homework » urban legend

### Simplex method

Dantzig's real contribution is not the definition of linear optimization but the invention of the **simplex method** to solve it

- Problem look continuous, but is in reality essentially combinatorial
- Optimal solution must lie on vertex of D
- Algebraic description of vertices + improving moves from vertex to vertex

### Algorithmic complexity

#### An efficient algorithm must

- Solve the problem (at least approximately)
- In a finite amount of time (must eventually stop)
- Or, better: in an amount of time bounded by a function of the problem size

#### **Essential distinction**

polynomial-time complexity (= easy)

– exponential-time complexity (= hard)

# Polynomial always beats exponential, eventually

n	$n^5$	$2^n$	$n^{5}/2^{n}$
20	3.05  seconds	1 second	3.05
40	1.7 minutes	12.1 days	$9.3 \times 10^{-5}$
60	12.9 minutes	349 centuries	$6.7 \times 10^{-10}$

General optimization is provable exponential Could linear optimization be polynomial?

# Algorithmic complexity for linear optimization methods

Georges Dantzig (1914-2005)

Proposes in 1947 the simplex method for linear optimization, very efficient in practice but whose algorithmic complexity is unknown

- V. Klee et G. Minty (1972)
  Show that the simplex method is exponential in the worst case
- L. Khachyan (1978)

Proves that the ellipsoid method (due to Shor) is polynomial in the worst; it is however very slow in practice!

• N. Karmarkar (1984)

Proposes a new type of method :interior-point methods, that are both polynomial in the worst case and very efficient in practice (often more than the simplex method)

### Optimization before Dantzig: a few pioneers

#### Simplex method is the start of modern era of optimization

- Leonid Kantorovich (1912-1986)
  Linear optimization en 1939! (Nobel prize in 1975 with Koopmans)
- John Von Neumann (1903-1957)
  Game theory (minimax theorem en 1928, link with linear optimiz.)
- Carl Friedrich Gauss (1777-1855)
   Steepest descent method (non-linear optimization)
- Leonhard Euler (1707-1783)
   Calculus of variations (infinite-dimensional optimization)
- Pierre de Fermat (1601-1665)
   Light: takes minimum-time path (already by Hero of Alexandria)
- Nature!
   In physics: least action principle, minimum energy principle

# An example of linear optimization

or

Can optimization help with your diet?

#### The diet problem

Eating not just a good breakfast, but an optimal breakfast.

Eight foods to choose from:

- 1. pancakes =  $x_1$
- 2. milk =  $x_2$
- 3. freshly squeezed orange juice =  $x_3$
- 4. scrambled eggs =  $x_4$
- 5. Weetabix, a high-fiber cereal =  $x_5$
- 6. fresh blueberries =  $x_6$
- 7. bacon =  $x_7$
- 8. vitamin/mineral pills =  $x_8$

Source: M. Wright



Source: M. Wright

Our goal is to find the cheapest breakfast satisfying daily fiber  $\geq 120$ , daily calcium  $\geq 100$ , and all x's  $\geq 0$ . (Can't eat -5 units of a food.)

food	fiber	calcium	cost
pancakes	5	10	3
milk	0	40	1
orange juice	4	5	5
eggs	8	0	3
Weetabix	100	20	10
blueberries	50	20	5
bacon	0	0	10
pills	30	400	2

Source: M. Wright

With this (apparently sensible) formulation, what's the optimal breakfast??



1.14 units of Weetabix and .19 units of vitamin/mineral pills.

Yuck. Not even milk for the Weetabix!

Source: M. Wright

### The original diet problem

Compute the optimal daily diet for US soldiers Satisfy nutritional requirements Minimize the cost

77 types of food, 9 nutritional requirements Solved in 1947 by Jack Laderman's team in 120 man-days

Optimal cost found: \$39.69 per year

# Duality

or

Can we trust optimization?

### Trust an optimal solution?

#### Difference between solving

- A system of equations :  $0 = f_1(x) = f_2(x) = ...$
- An optimization problem: min f(x)

#### How?

- For a system of equations: simply replace x in  $f_1(x)$ ,  $f_2(x)$ , ...
- For an optimization problem : how ?

#### Generate bounds for linear optimization

Consider the small problem

$$f^* = \max x_1 + 2x_2 + 3x_3$$
 such that  $\begin{cases} x_1 + x_2 \le 1 & (a) \\ x_2 + x_3 \le 2 & (b) \\ x_2 + 3x_3 \le 9 & (c) \\ x_3 \le 3 & (d) \end{cases}$ 

- a. Vector x=(1,0,2) is feasible and corresponds to objective value equal to 7
  - $\Rightarrow$  lower bound  $f^* > 7$
- b. Combine the constraints according to (a) + (c)  $(x_1 + x_2) + (x_2 + 3x_3) \le 1 + 9 \Leftrightarrow x_1 + 2x_2 + 3x_3 \le 10$  $\Rightarrow$  upper bound  $f^* \le 10$

#### Finding better lower bounds

$$f^* = \max x_1 + 2x_2 + 3x_3$$
 such that  $\begin{cases} x_1 + x_2 \le 1 & (a) \\ x_2 + x_3 \le 2 & (b) \\ x_2 + 3x_3 \le 9 & (c) \\ x_3 \le 3 & (d) \end{cases}$ 

a. Lower bound can be improved: find a better vector x = (2, -1, 3) gives an objective function equal to 9  $\Rightarrow$  lower bound  $f^* \geq 9$ 

To find better vectors we may use the simplex method

#### Finding better upper bounds

$$f^* = \max x_1 + 2x_2 + 3x_3 \text{ such that } \begin{cases} x_1 + x_2 \le 1 & (a) \\ x_2 + x_3 \le 2 & (b) \\ x_2 + 3x_3 \le 9 & (c) \\ x_3 \le 3 & (d) \end{cases}$$

- a.  $x = (2, -1, 3) \Rightarrow \text{lower bound } f^* \geq 9$
- b. Upper bounds can also be improved: combining the constraints according to (a)+(b)+2(d) $x_1+x_2+x_3+2x_3 \le 1+2+2\times 3 \Leftrightarrow x_1+2x_2+3x_3 \le 9$  $\Rightarrow$  upper bound  $f^* \le 9 \Rightarrow (2,-1,3)$  is optimal!

#### How to find better upper bounds?

$$f^* = \max x_1 + 2x_2 + 3x_3$$
 such that  $\begin{cases} x_1 + x_2 \le 1 & (a) \\ x_2 + x_3 \le 2 & (b) \\ x_2 + 3x_3 \le 9 & (c) \\ x_3 \le 3 & (d) \end{cases}$ 

Finding upper bound is an optimization problem!

Combining the constraints as  $y_1(a) + y_2(b) + y_3(c) + y_4(d)$  we should

- $\diamond$  choose nonnegative  $y_1, y_2, y_3, y_4$  nonnegative
- $\diamond$  satisfying  $y_1 = 1, y_1 + y_2 + y_3 = 2, y_2 + 3y_3 + y_4 = 3$
- $\diamond$  minimizing value of the bound  $y_1 + 2y_2 + 9y_3 + 3y_4$

#### The dual problem

$$f^* = \max x_1 + 2x_2 + 3x_3$$
 such that  $\begin{cases} x_1 + x_2 \le 1 & (a) \\ x_2 + x_3 \le 2 & (b) \\ x_2 + 3x_3 \le 9 & (c) \\ x_3 \le 3 & (d) \end{cases}$ 

and

$$d^* = \min y_1 + 2y_2 + 9y_3 + 3y_4 \text{ such that } \begin{cases} y_1 + y_2 + y_3 = 2 \\ y_2 + 3y_3 + y_4 = 3 \\ y_1, y_2, y_3, y_4 \ge 0 \end{cases}$$

form a pair of dual problems!

#### Duality properties

#### a. Weak duality

Any feasible vector for the dual problem provides an upper bound for all feasible vectors of the problem (and vice-versa!)

#### b. Strong duality

There *always* exists a feasible dual vector that certifies optimality of the primal solutions

Therefore both problems have the *same* optimal value!

This optimal dual vector is an optimality certificate, can be found by any means (guessing, simplex method, etc.) and can be checked easily and independently

# Duality summarized

- Duality is a way to certify quality or optimality of some feasible vectors
- Computing dual certificate is itself a linear optimization problem
- Some methods automatically compute the dual optimal solution (such as the simplex method)
- Dual variables have an economic interpretation (they act as prices for the constraints)

# Performance

Or

Can Optimization compete with computers?

# Progress for linear optimization

In 1947, Laderman solves a small diet problem (77 ingredients x 9 nutriments) in 120 (wo)men-days (with Dantzig's simplex algorithm)

In 2008, linear problems with millions of variables/constraints are routinely solved on standard desktop computer (provided enough memory is available)

Is it really remarkable?

Modern computers are so powerful ...
Is this a simple consequence of Moore's law, i.e. speed doubles every 1.5 year?

### A 25-year comparison: 1987-2002

- R. Bixby considers, for the same problem
- The 1987 version of his solver CPLEX on a 1987 computer: 2.5 hours (~10 000 s)
- The same 1987 version CPLEX on a 2002 cmputer : 10 seconds

This matches Moore's law quite well

### A 25-year comparison: 1987-2002

- R. Bixby also compares on a larger problem:
- The 1987 version of CPLEX on a 2002 computer: 8 hours (~30 000 s)
- The 2002 version of CPLEX on a 2002 computer: 30 seconds

Algorithmic progress matches technological progress: 1000x each

Total gain: 6 orders of magnitude in 15 years

# Summary

# Optimization is exciting

- Can be used everywhere
- Is actually used everywhere!

Provably difficult ... but we can manage

- Many large-scale problems can be solved efficiently in practice
- A lot of research still going on

### There is a lot more to say ...

- Nonlinear Optimization (and convexity)
- Discrete Optimization
- Convex/Structured optimization
- Combinatorial Optimization
- Global Optimization
- Heuristic and meta-heuristic methods
- Average case analysis
- etc.

Thank you for your attention

# Can you match an optimization solver?

